Lazy Segment Trees

CS 491 - Competitive Programming

By Andrea Zhou

Last Lecture: Segment Trees

Objectives:

- Build a Segment Tree
- Query a Segment Tree

Consider the following data array

[85, 61, 75, 59, 49, 64, 50, 37]

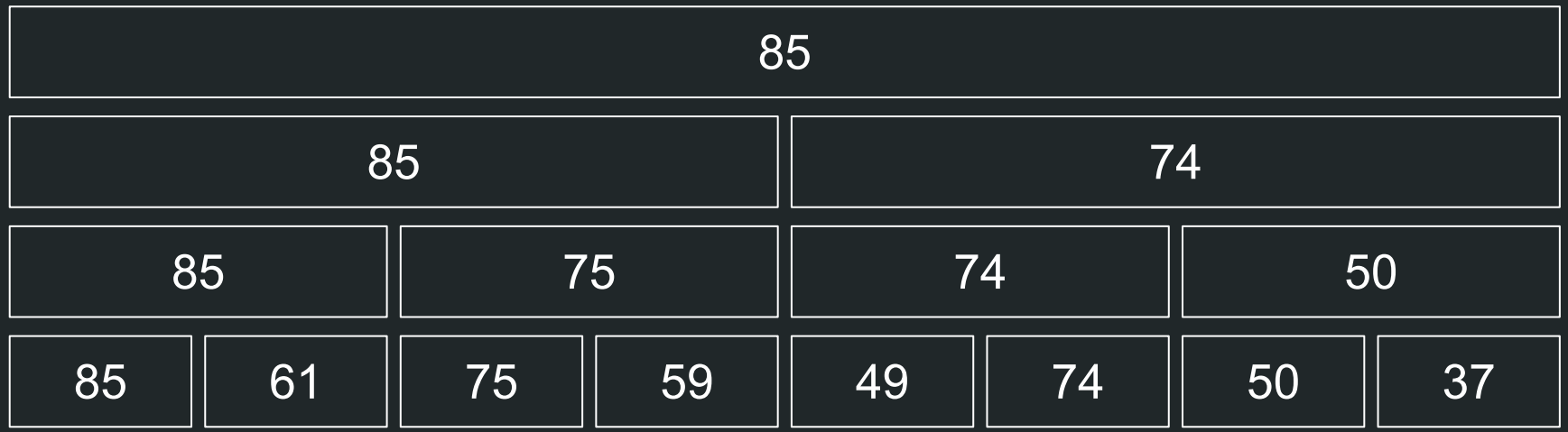
Consider the following data array

[85, 61, 75, 59, 49, 64, 50, 37]

If we want to perform range queries such as

- Find the maximum element in the range [5,6]
- Find the maximum element in the range [0,4]
- Find the maximum element in the range [3,7]
- ...

Segment Tree



But what if we wanted to change some of the data?

- Add 5 to elements in the range [4,7]
- Find the maximum element in the range [5,6]
- Find the maximum element in the range [0,4]
- Subtract 2 from elements in the range [6,6]
- Find the maximum element in the range [3,7]
- ...

But what if we wanted to change some of the data?

- Add 5 to elements in the range [4,7]
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- Subtract 2 from elements in the range [6,6]
- Find the maximum element in the range [3,7]
- ...

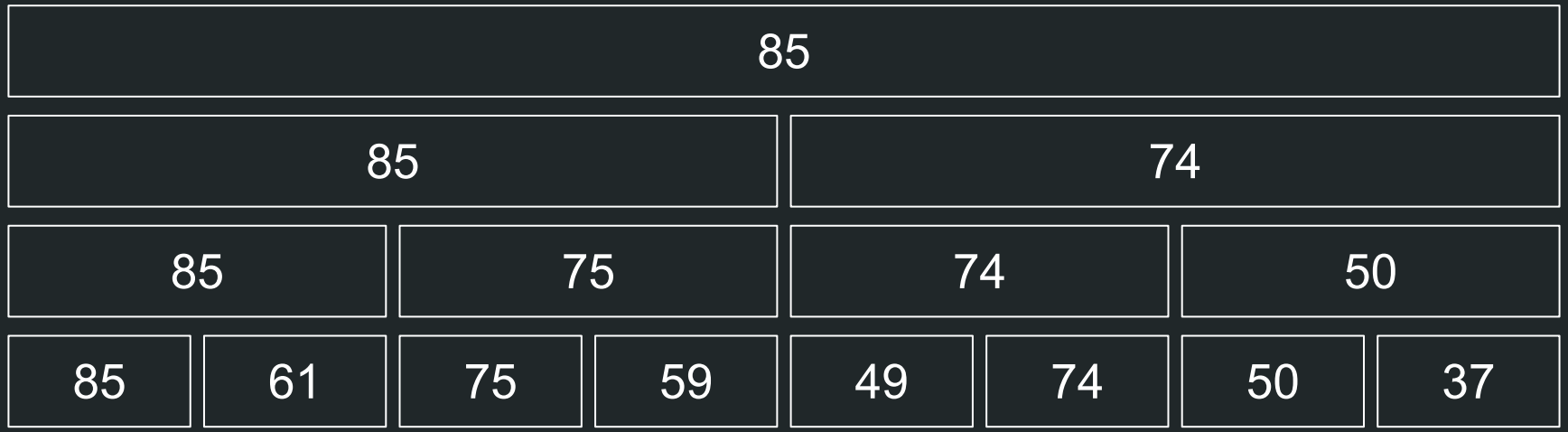
← might not return correct max without updating ST

Today’s Lecture: Lazy Propagation in Segment Trees

Objectives:

- Perform updates to data while still being able to query the segment tree in log(n) time

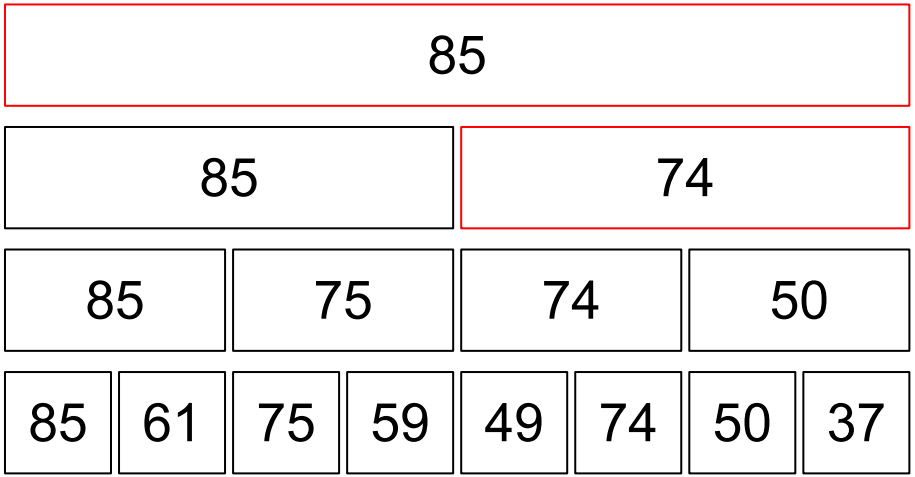
Segment Tree



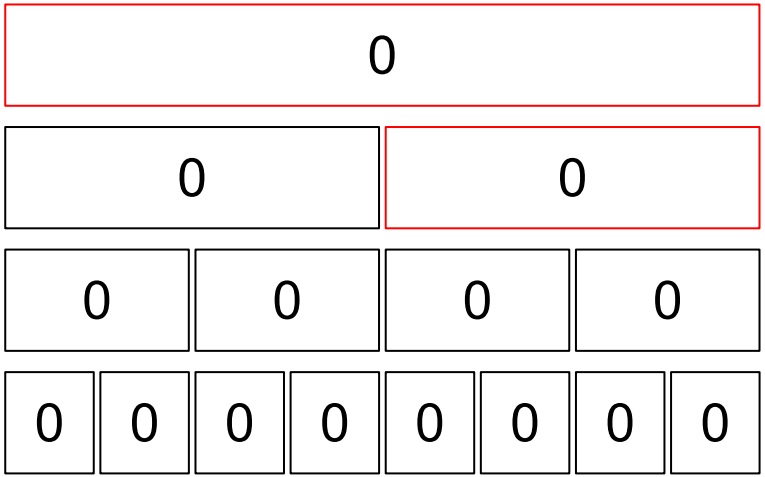
1. Update: Add 5 to elements in the range [4,7]
2. Query: Find the maximum element in the range [5,6]
3. ...

1. “Add 5 to elements in the range [4,7]”

Segment Tree



Lazy Tree



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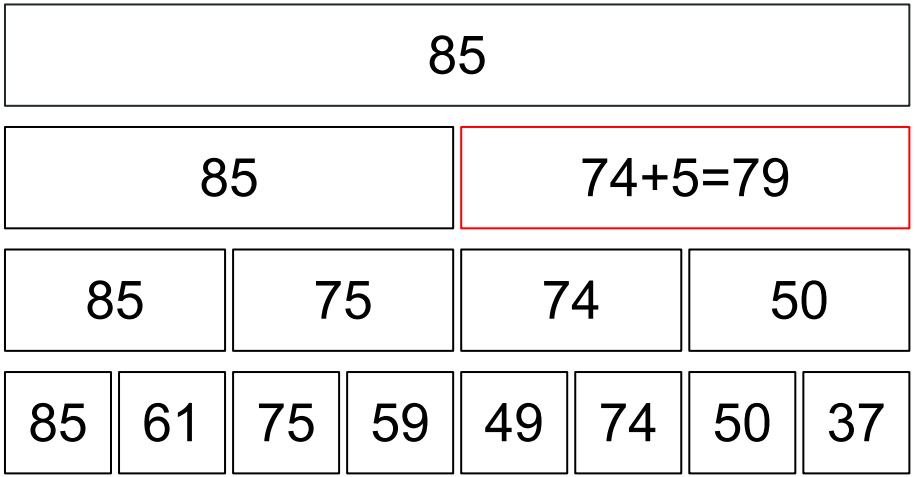

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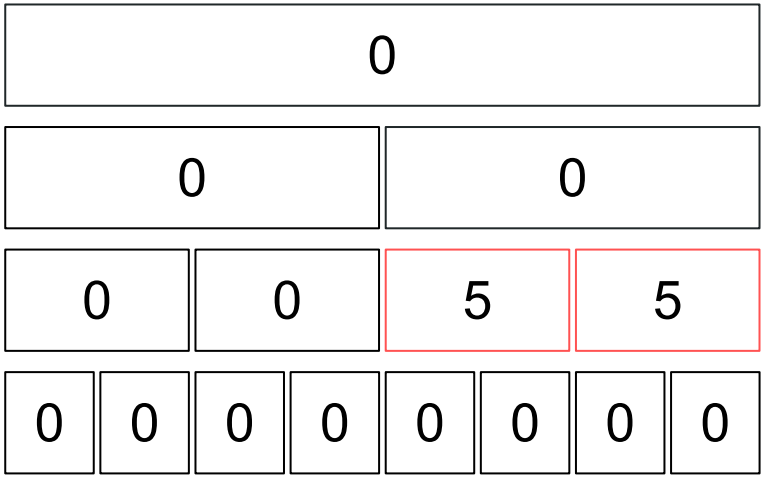

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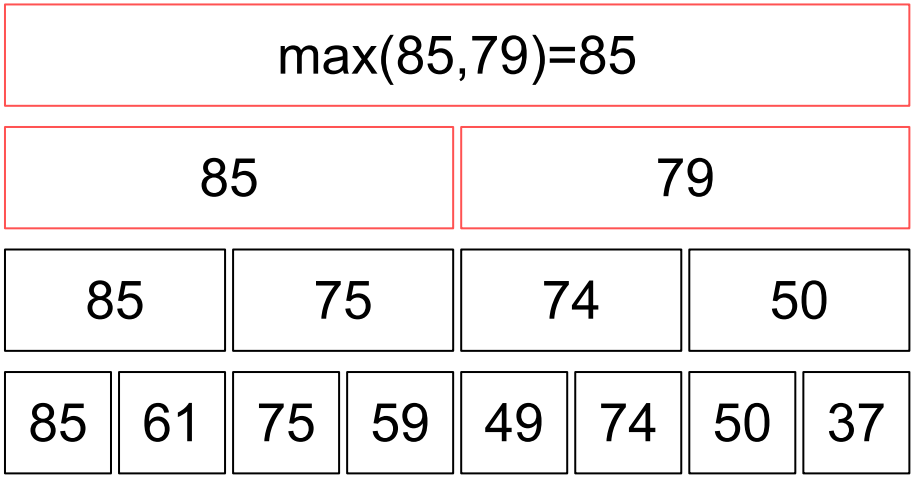

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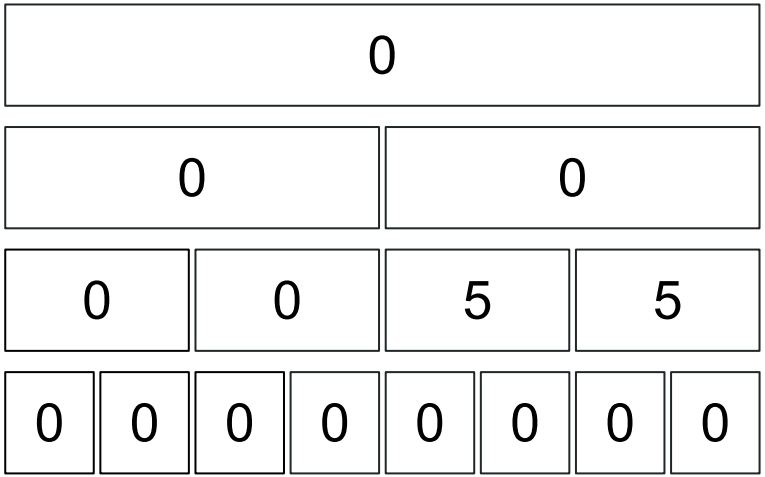

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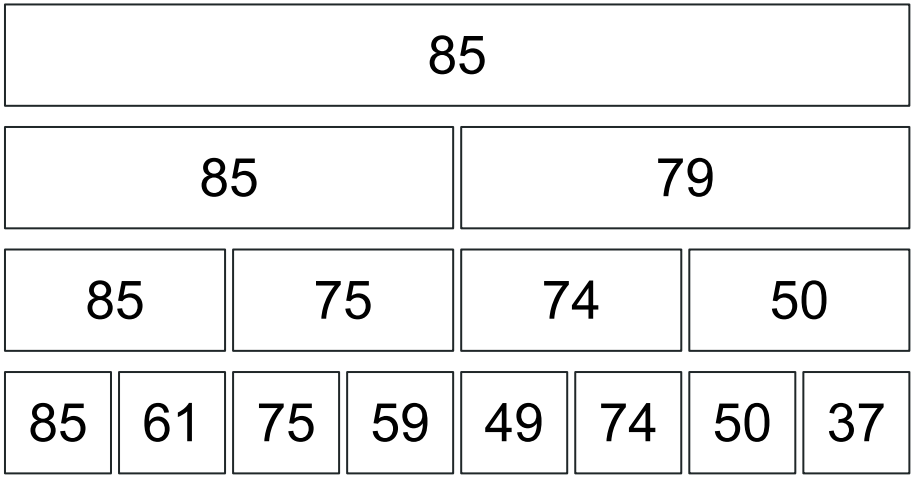

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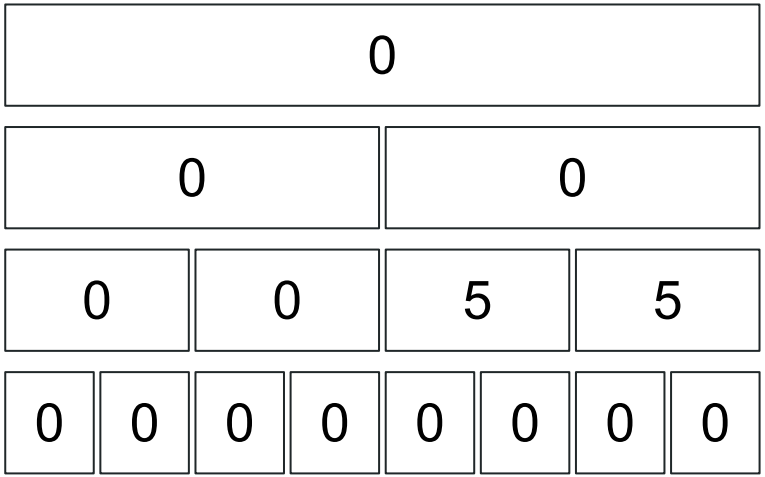

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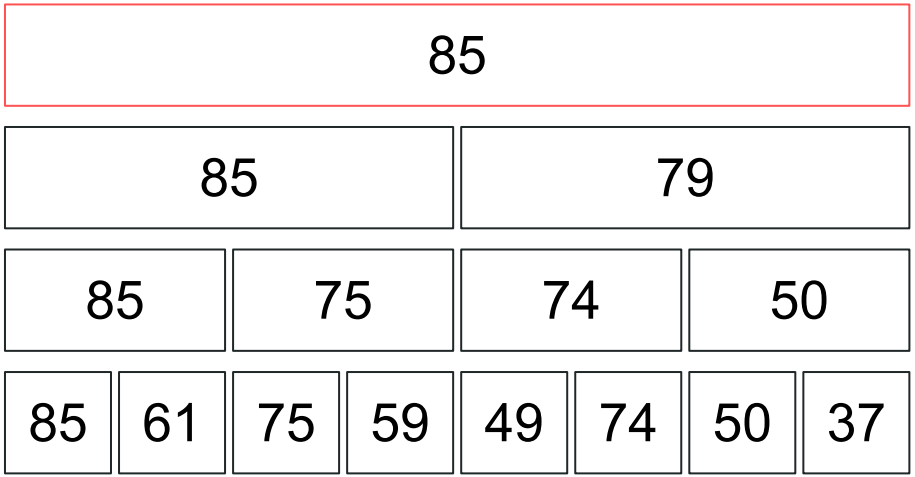

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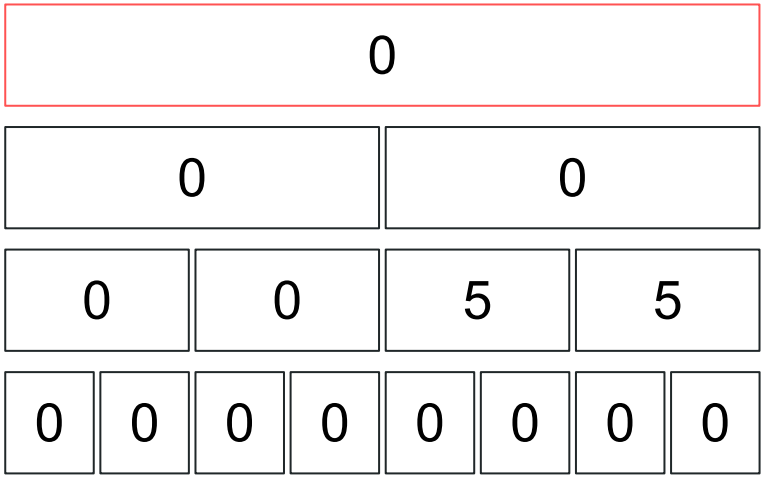

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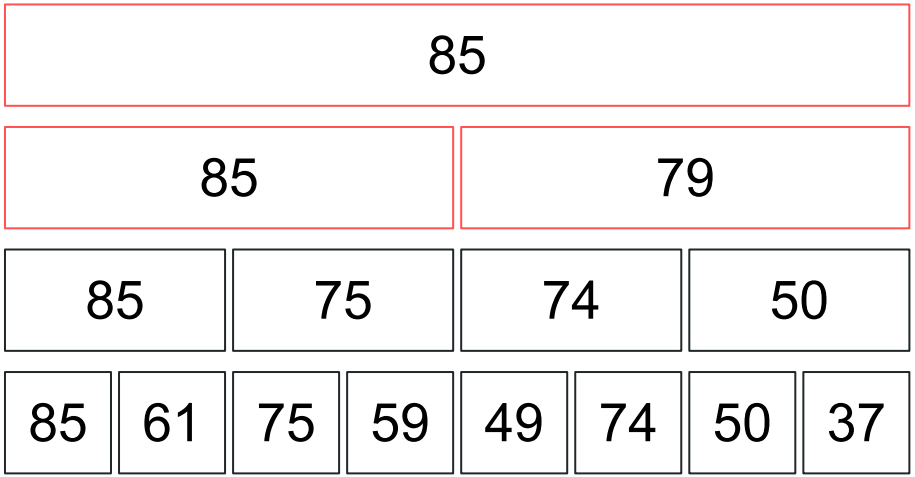

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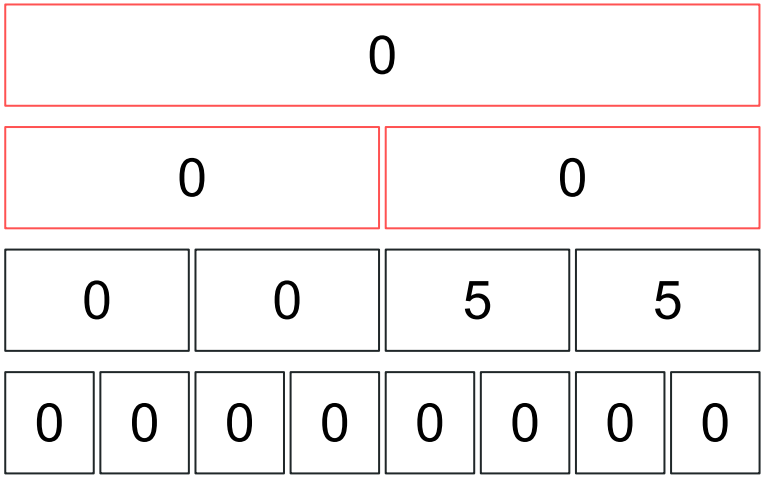

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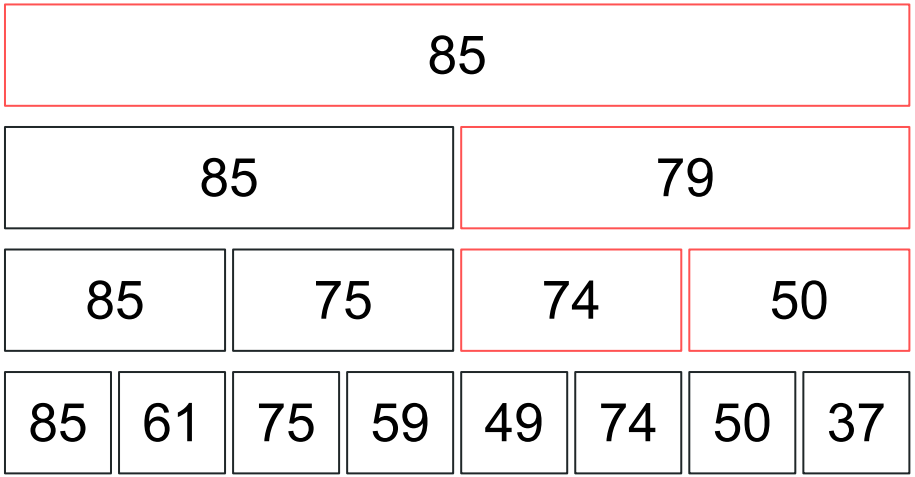

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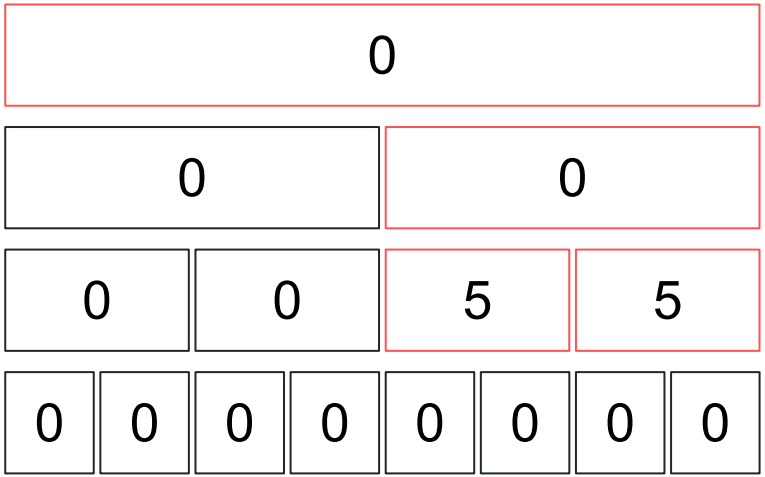

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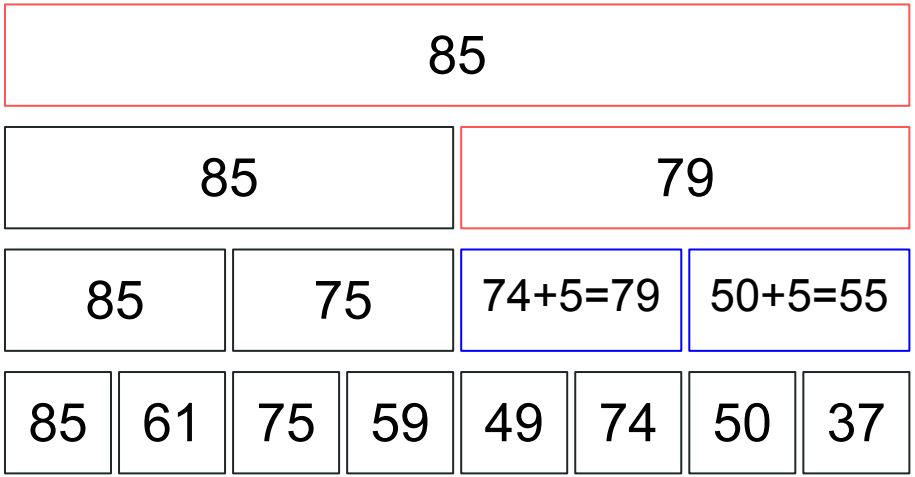

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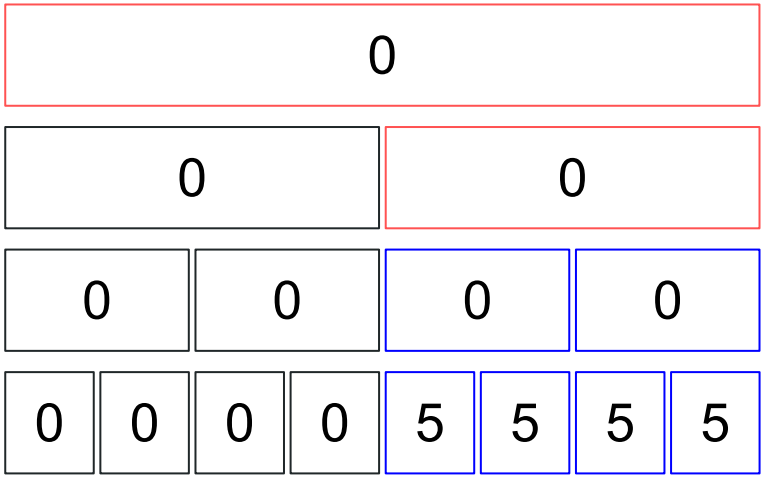

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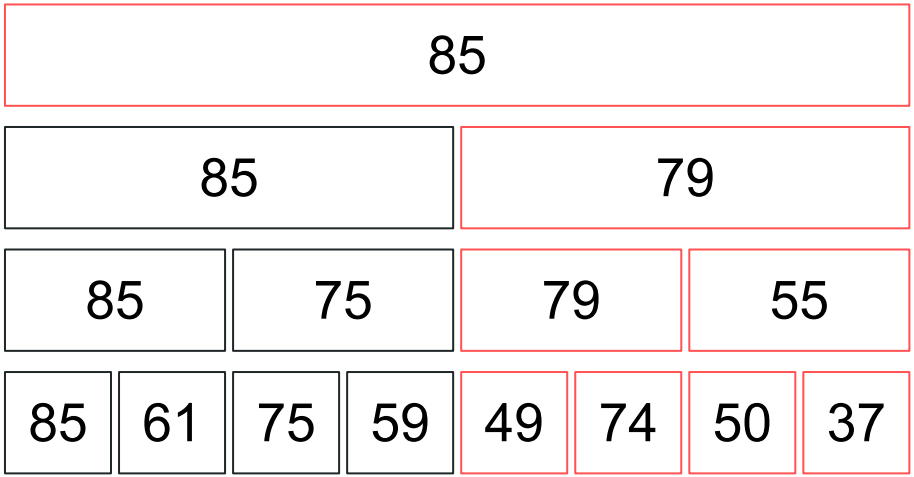

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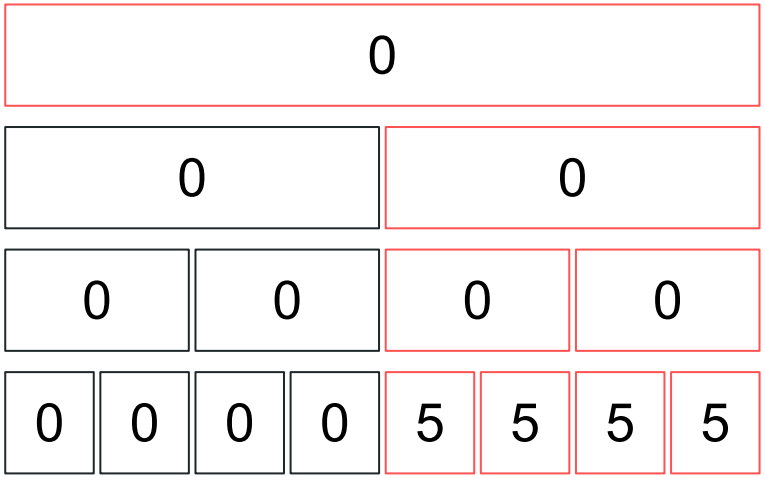

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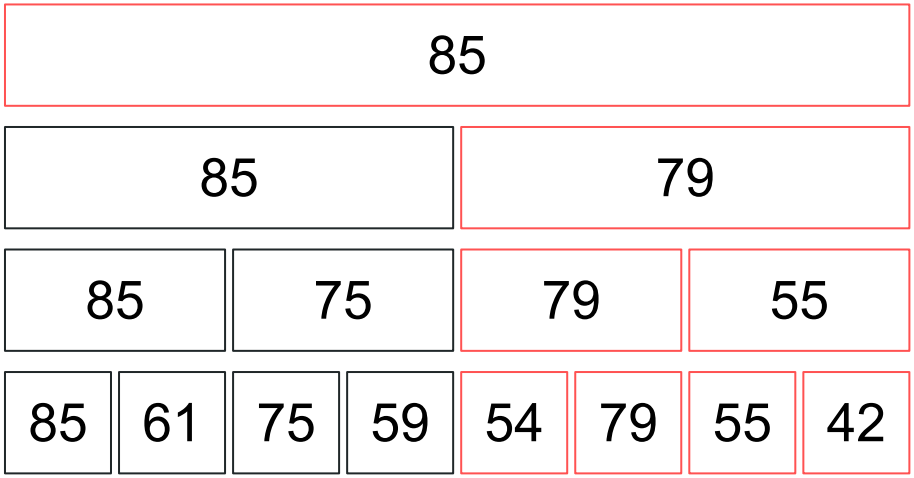

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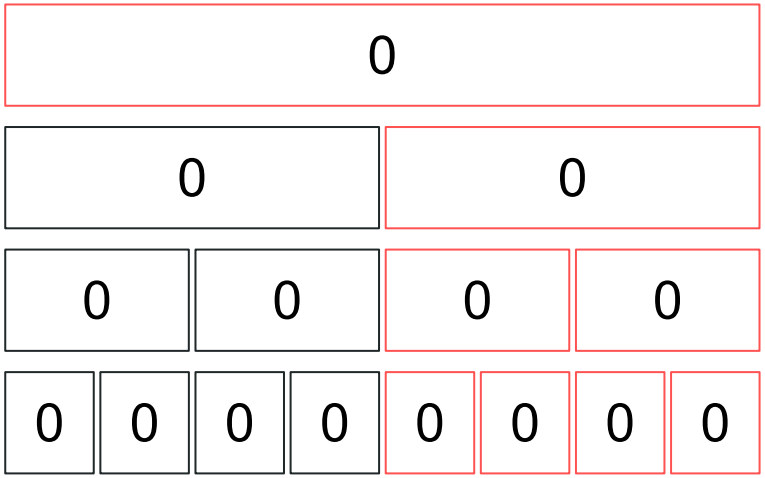

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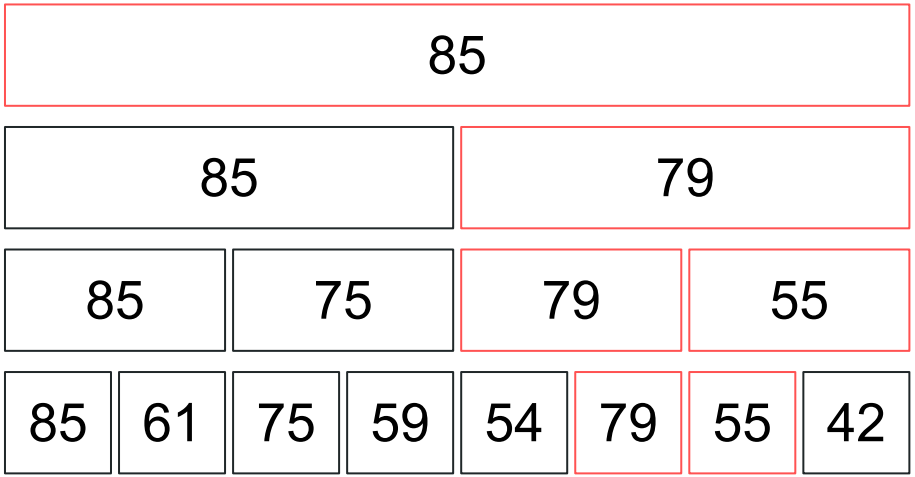

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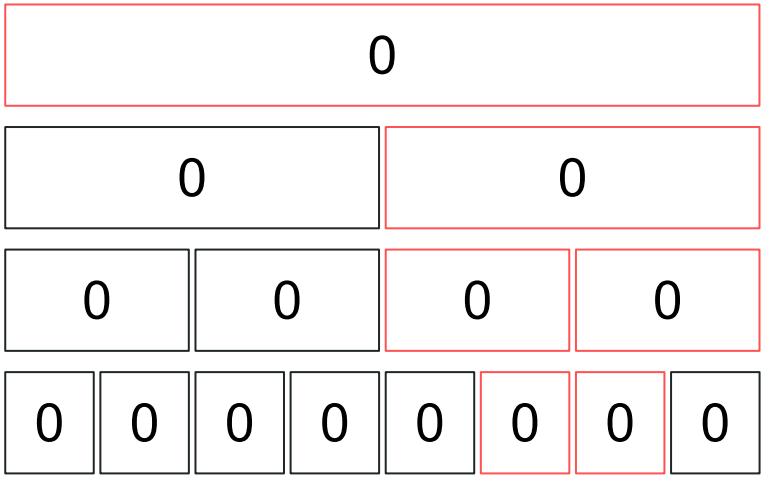

CP4 Segment Tree Class

```
1  class SegmentTree {
2     private:
3        int n; // size
4        vi A, st, lazy;
5        int l(int p) { return p<<1; } // go left
6        int r(int p) { return (p<<1)+1; } // go right
7
8        int conquer(int a, int b) {
9          if (a == -1) return b; // corner case
10         if (b == -1) return a;
11         return min(a, b); // RMQ
12       }
```

Code provided from Dr. Mattox Beckman

Building the Segment Tree

```
13  void build(int p, int L, int R) { // O(n)
14    if (L == R) st[p] = A[L]; // base case
15    else {
16      int m = (L+R)/2;
17      build(l(p), L , m);
18      build(r(p), m+1, R);
19      st[p] = conquer(st[l(p)], st[r(p)]);
20    } }
```

Code provided from Dr. Mattox Beckman

Querying

- L and R give you the bounds with respect to the orignal array.
- i and j give you the bounds for the query

```
21 int RMQ(int p, int L, int R, int i, int j) { // O(log n)
22   propagate(p, L, R); // lazy propagation
23   if (i > j) return -1; // infeasible
24   if ((L >= i) && (R <= j)) return st[p]; // found the seg
25   int m = (L+R)/2;
26   return conquer(RMQ(l(p), L ,  m, i , min(m, j)),
27                  RMQ(r(p), m+1, R, max(i, m+1), j ));
28 }
```

Code provided from Dr. Mattox Beckman

Updating

```
29  void update(int p, int L, int R, int i, int j, int val) {
30    propagate(p, L, R); // lazy propagation
31    if (i > j) return;
32    if ((L >= i) && (R <= j)) { // found the segment
33      lazy[p] = val; // update this
34      propagate(p, L, R); // lazy propagation
35    } else {
36      int m = (L+R)/2;
37      update(l(p), L , m, i , min(m, j), val);
38      update(r(p), m+1, R, max(i, m+1), j , val);
39      int lsub = (lazy[l(p)] != -1) ? lazy[l(p)] : st[l(p)];
40      int rsub = (lazy[r(p)] != -1) ? lazy[r(p)] : st[r(p)];
41      st[p] = (lsub <= rsub) ? st[l(p)] : st[r(p)]; } }
```

Code provided from Dr. Mattox Beckman

Propagating

```
42  void propagate(int p, int L, int R) {
43    if (lazy[p] != -1) { // has a lazy flag
44      st[p] = lazy[p]; // [L..R] has same value
45      if (L != R) // not a leaf
46        lazy[l(p)] = lazy[r(p)] = lazy[p]; // propagate
47      else // L == R, a single index
48        A[L] = lazy[p]; // time to update this
49      lazy[p] = -1; // erase lazy flag
50    } }
```

Code provided from Dr. Mattox Beckman